

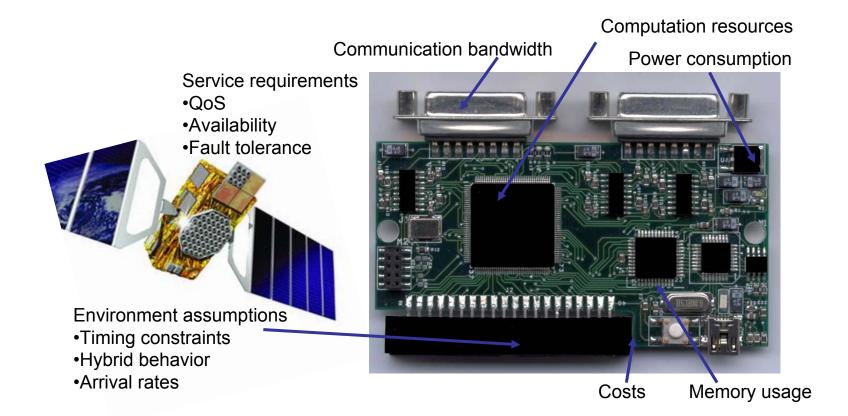
Quasimodo

Quantitative System Properties in Model-Driven-Design of Embedded Systems

Kim G. Larsen & Brian Nielsen Aalborg University, DK

Quantitative System Properties

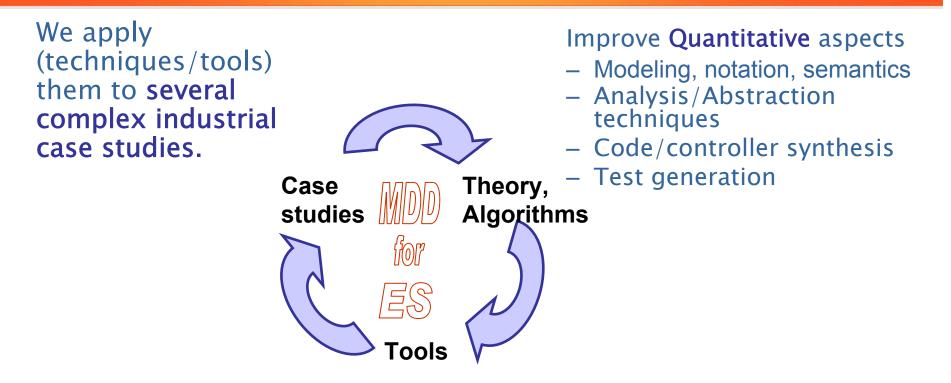




Page 2

Quasimodo Research Goals



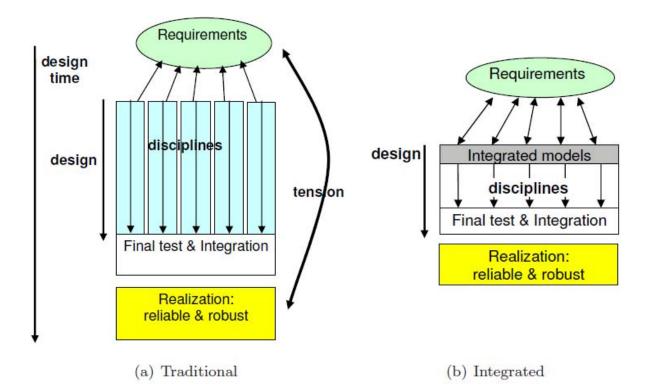


- Provide unique tool components to be used as plug-ins in industrial tools / tool chains
- Create first prototypes of an integrated tool environment
- Disseminate

Page 3

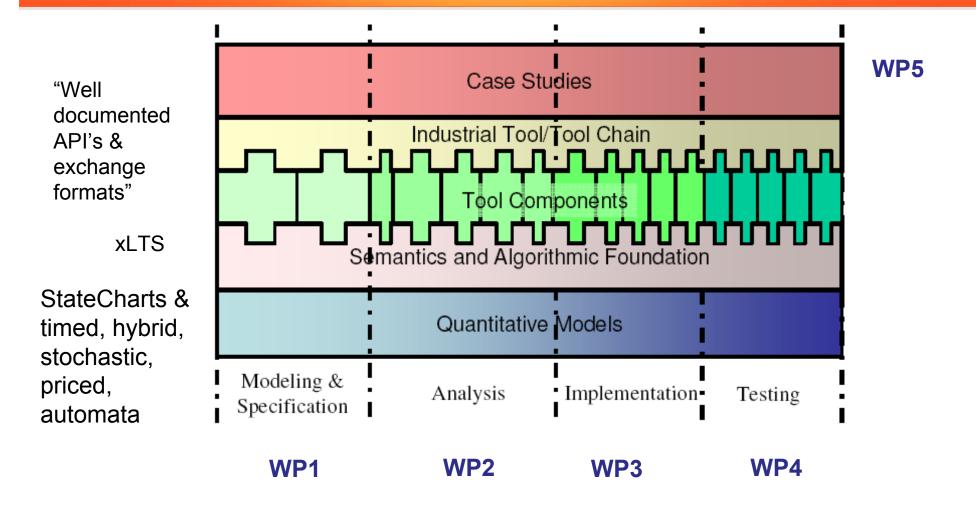
Early Integration & validation





Improve competitiveness of European companies that rely on the design and integration of embedded systems in their products by reducing costs and time to market





Partners asimodo •Aalborg University / CISS •Terma Space A/S 2 ötebo •Embedded Systems Glasgow Edinburgh Institute North Öland DENMARK Copenhagen *• Malmö UNITED Sea -Twente University –Radboud University Man (U.K.) Leeds •Saarland University Manchester AND •CHESS iverpoo KINGDOM •RWTH Aachen Amsterdam NETH. Poznań Birmingham ardiff •Hydac Electronic Gmbh Rotterdam London ★ Leipzig •Université Libre de Bonn GERMANY

Liffe

FRANCE

Toulouse

Andorra

Marso

Guernsey (U.K.)

Jersey (U.K.)

Bilbao

Nantes

Bordeaux

Bruxelles (CFV)

•CNRS (LSV, ENS Cachan)

BEL.

Strasbou

MONACI

Frankfurt am Main

Stuttgart

Munich

MARINO

Florence

Luxembourg

CZECH REPUBL

AUSTRIA

CROATIA



Work Tasks



WP1 Modeling and Specification	WP2 Analysis	WP3 Implementation	WP4 Testing	WP5 Case Studies, Tools, Disse- mination and Exploitation
T1.1 Model Process Improvement	T2.1: State Space Representation and Model Checking	T3.1: Controller synthesis and scheduling	T4.1: Test Generation	T5.1: Case Studies
T1.2: Modeling of Quantitative System Aspects	T2.2: Abstraction, Refinement and Compositionality	T3.2: Implementability and code generation	T4.2: Approximate Testing	T5.2: Tool Plugins and Tool Chain Integration.
T1.3: Design Notation and Tools	T2.3: Approximate Analysis Techniques			T5.3: Dissemination and Exploitation







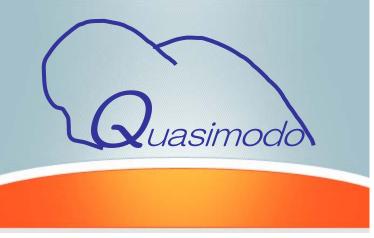
www.quasimodo.aau.dk







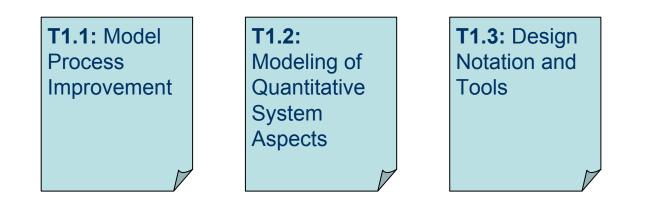


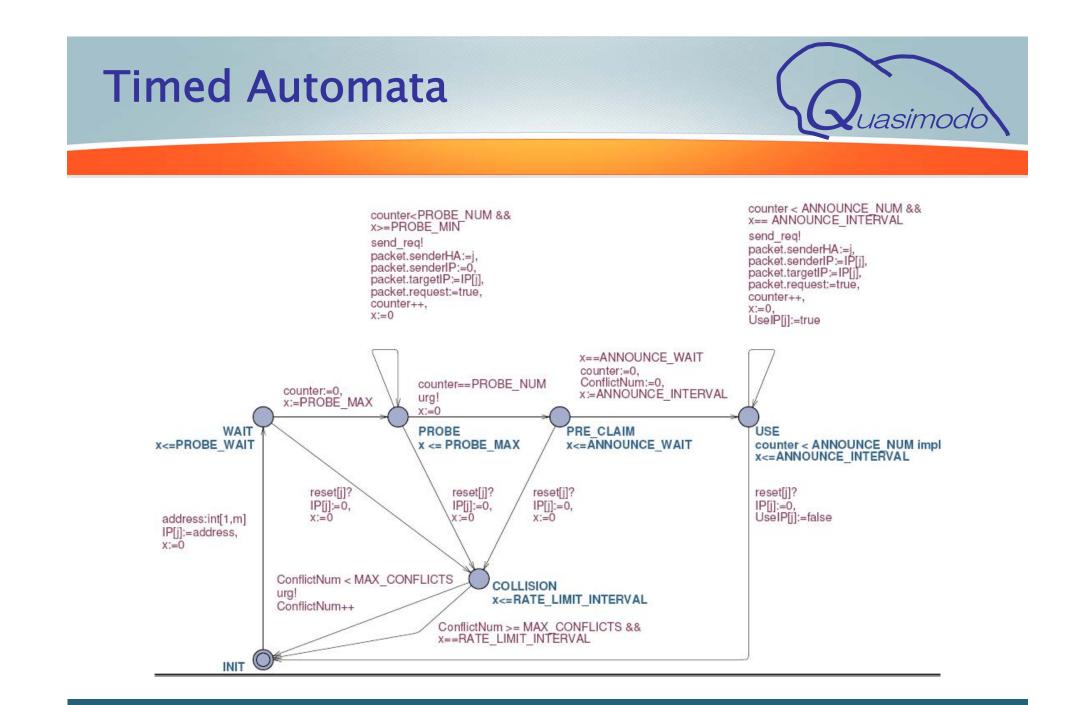


WP1: Modeling and Specification $\zeta_{($

Leader: Frits Vaandrager (ESI/RU)

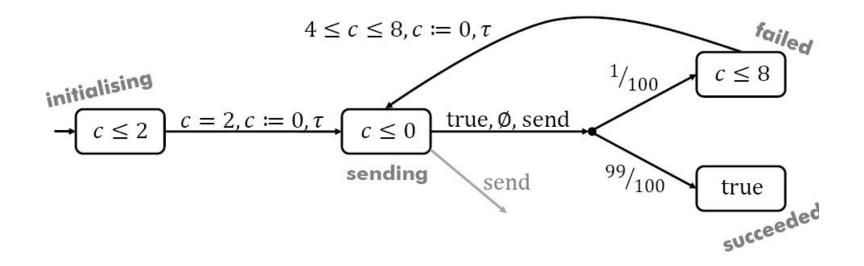
- How to obtain adequate and faithful models of embedded systems?
 - Modeling formalisms integrating timed, hybrid, cost and stochastic aspects
 - Overarching design notation for embedded systems with accompanying tool support?
 - Modeling process and model management



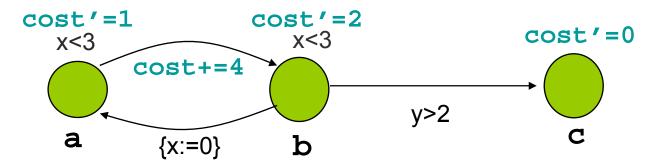




Timed automata + probabilistic branching







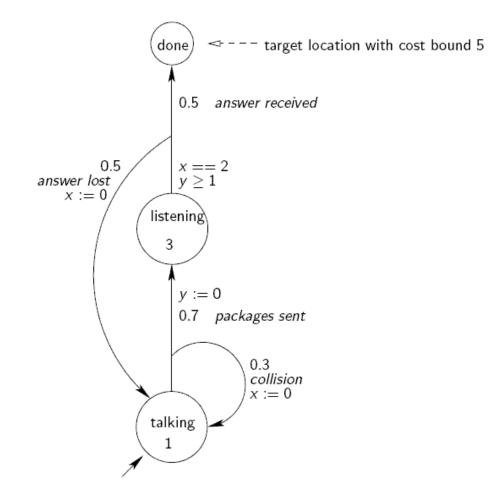
Timed Automata + Costs on transitions and locations
 Cost of performing transition: Transition cost
 Cost of performing delay d: (d x location cost)

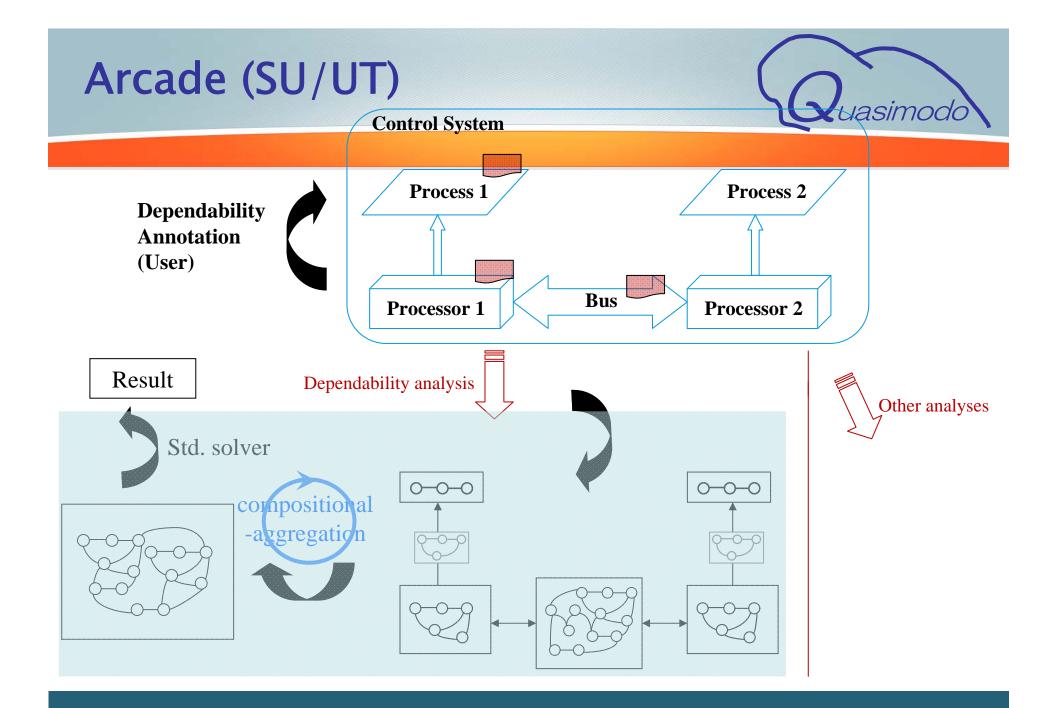
$$(\mathbf{a}, \mathbf{x}=\mathbf{y}=0) \xrightarrow{\mathbf{a}} (\mathbf{b}, \mathbf{x}=\mathbf{y}=0) \xrightarrow{\frac{\varepsilon(2.5)}{\mathbf{2.5 \times 2}}} (\mathbf{b}, \mathbf{x}=\mathbf{y}=2.5) \xrightarrow{\mathbf{0}} (\mathbf{a}, \mathbf{x}=0, \mathbf{y}=2.5)$$

Cost of Execution Trace: Sum of costs: 4 + 5 + 0 = 9
 Problem: Find minimum cost of reaching location c

Priced Probabilistic Timed Automata





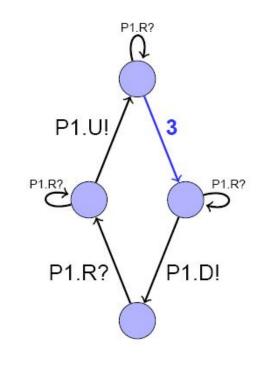


I/O Interactive Markov Chains



Input/Output Interactive Markov Chains

- Combination of I/O automata and Markov chains, close to IMCs
- Discrete state space
- Labelled transitions
 - Markovian transitions (3)
 - Input transitions (?)
 - Output transitions (!)
 - Internal transitions (;)
- Input enabled, outputs cannot be delayed
- Parallel composition, hiding, renaming, lumping, etc.



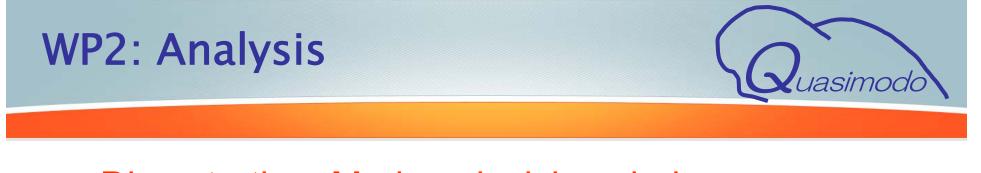
WP2: Analysis



Leader: Joost-Pieter Katoen (RWTH)

- How to accurately, effective and efficiently (algorithmically) analyze quantitative models?
 - Design of data structures for representing and exploring behavior of (combined) quantitative models.
 - Support for interrelating different quantitative models preserving properties.
 - Methods for allowing partial analysis of very complex models (size or quant. aspects considered).





- Discrete-time Markov decision chains:
 - Trace equivalence for labeled MDPs
 - Decision algorithms for (strong and weak) probabilistic simulation on DTMCs
 - Probabilistic CEGAR for MDPs
 - Regular expressions for PCTL counterexamples
- Continuous-time Markov chains:
 - Discrete-event simulation for CSL model checking
 - Parameter synthesis for time-bounded reachability
 - Infinite-state CTMC model checking
 - Advances in three-valued abstraction
 - Minimization of acyclic phase-type distributions



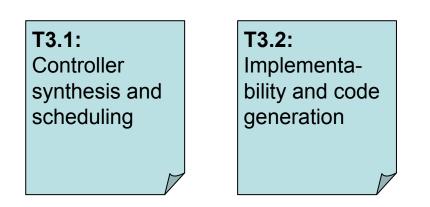
- Probabilistic timed automata:
 - Complexity results for 1-2C PTA model checking
 - Probabilistic semantics of timed automata
 - Probabilistic time-abstract bisimulation
 - Uppaal-PRO tool for PTA model checking
 - Undecidability of CBPR in priced 3C PTA
 - Decidability and data structures for Pareto-optimal reachability in multi-priced TA
 - Decidability for optimal infinite scheduling for priced TA using mean pay-offs and discounting metrics
 - Heuristic guided search for TA using Russian Doll principle

WP3: Implementation

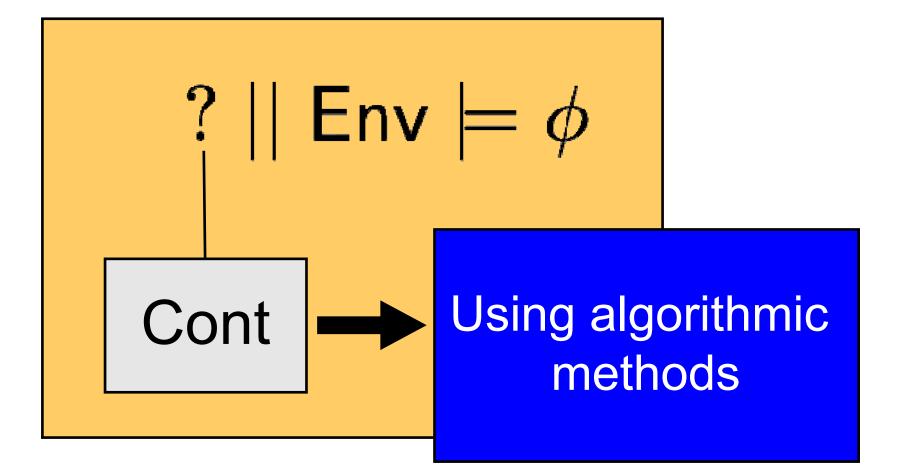


Leader: Jean Francois Raskin (CFV/ULB)

- How to algorithmically synthesize correct implementations of (continuous time) models on imperfect (digital) hardware/ software components?
 - Controller and scheduler synthesis
 - Property preserving code generation

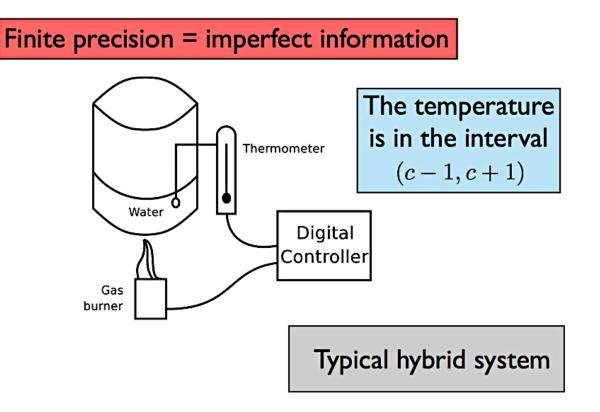


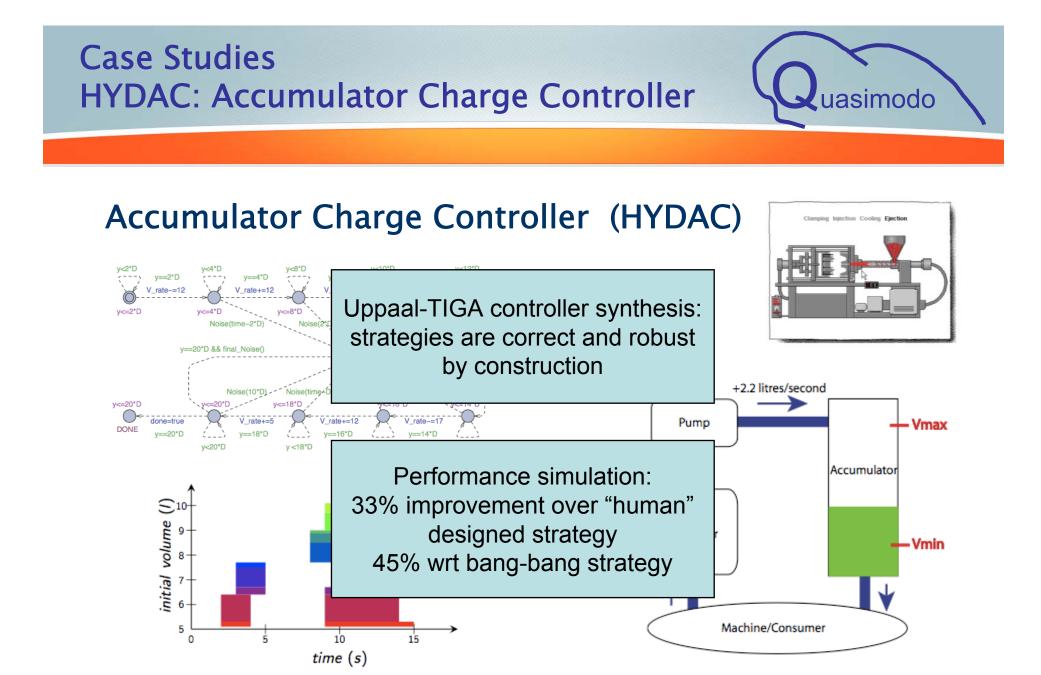






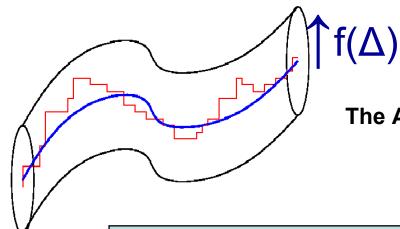
Highlight: imperfect information





Implementability and code generation





The AASAP semantics define a tube a strategies

- AASAP semantics define a "tube" of strategies instead of a unique strategy in the ASAP semantics.
- This tube can be refined into an implementation while preserving all LTL properties

Highlight: robustness analysis in UppAal



- Robust semantics is available within UppAal;
- Original results were formulated on regions;
- For implementation, we need a zone based formulation;

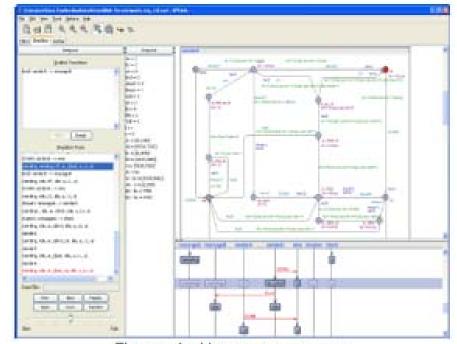


Figure 1: UPPAAL on screen.

WP4: Testing

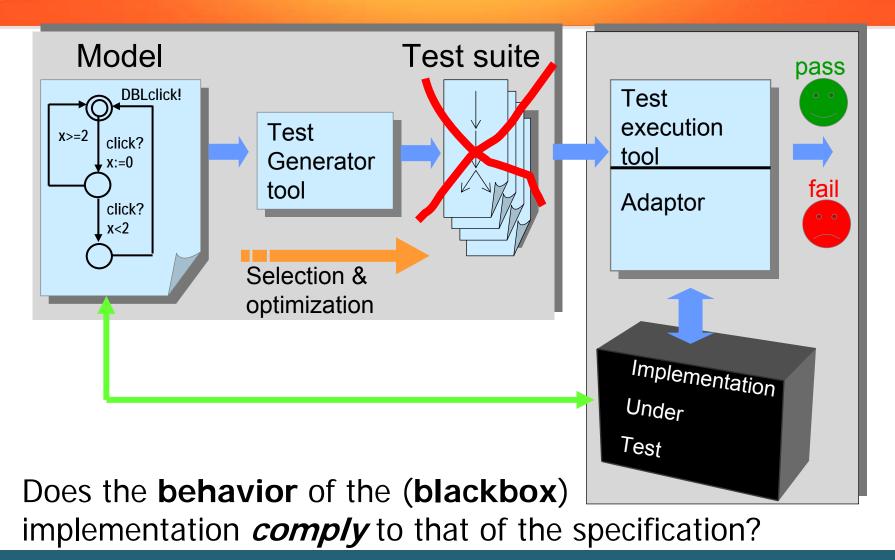


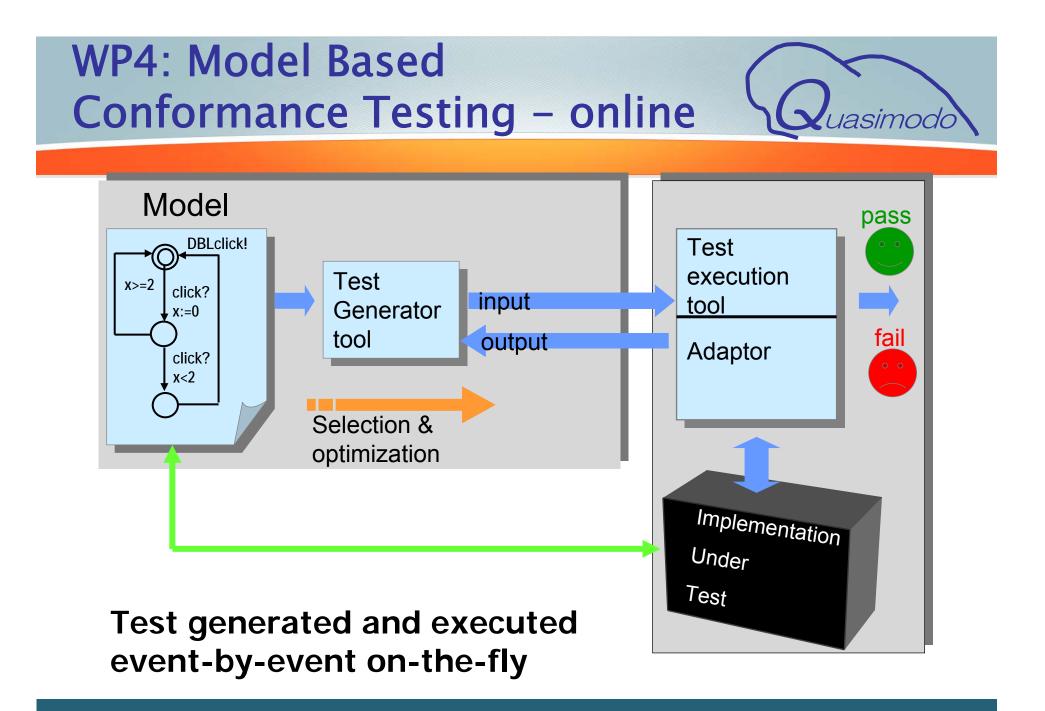
Leader: Arne Skou (AAU)

- How to automatically generate test cases from models with quantities?
 - Theory, algorithms, tools
 - Test selection methods, guidance and heuristics for optimal testing, and to define coverage metrics
 - Approximative testing methods to handle measurement and observation imprecisions

T4.1: Test generation	T4.2: Approximate Testing

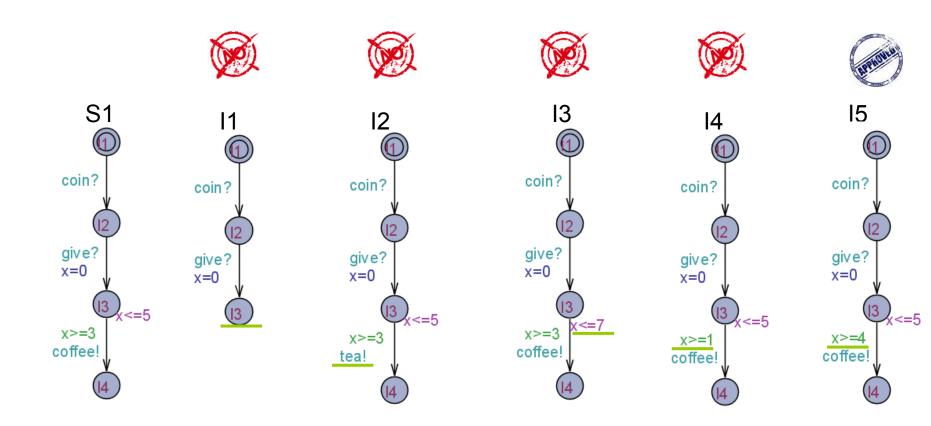
WP4: Model Based Conformance Testing – offline Quasimodo





Does In (rt-ioco) conform-to S₁ ?

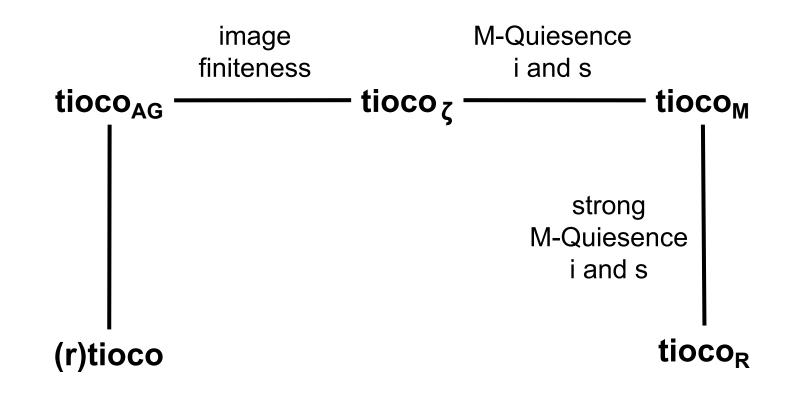






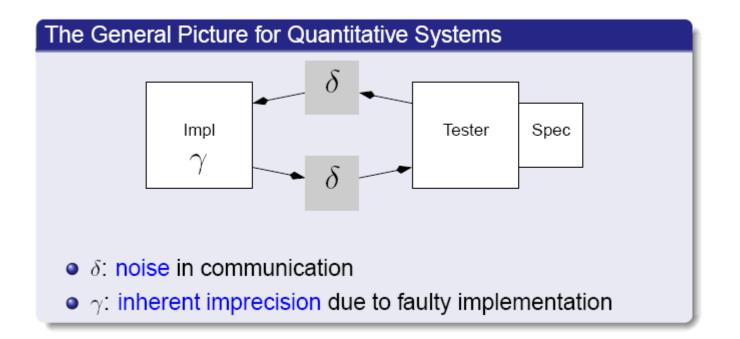
- Further work on timed conformance:
 - Conformance wrt. quiescence:
 - tioco_ζ: Unbounded delays are observable
 - tioco_M: Delays up to M are observable
 - tioco_R: Delays up to M and tau-moves are observable





Quantitative Testing Theory

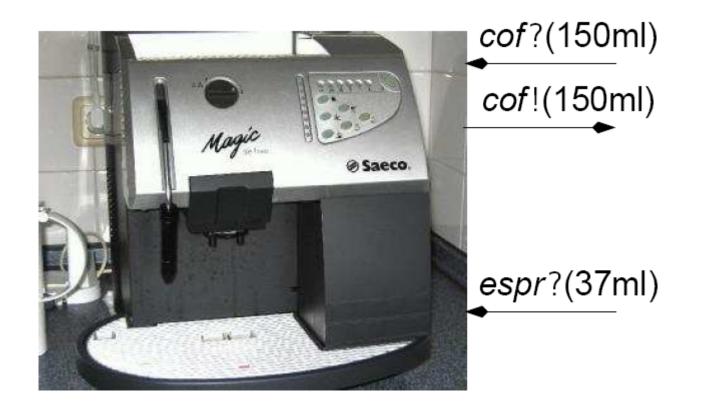




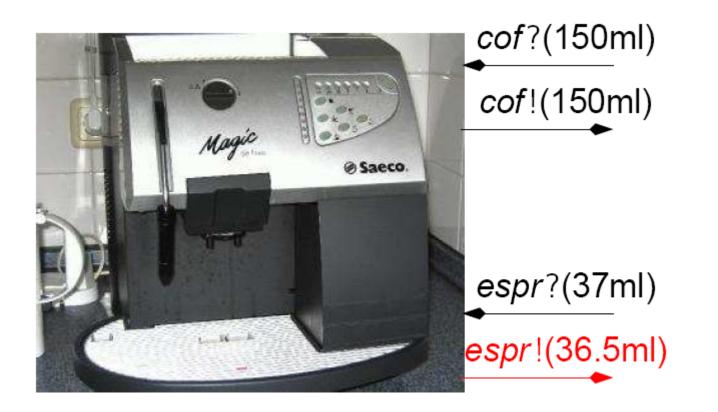
Results:

- A new conformance relation defining the 'distance' from Spec
- Sketch of algorithms for testing the 'distance'









Quasimodo Review, Bruxelles, February 19, 2009

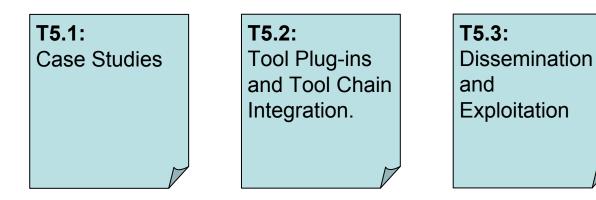
Page 35

WP5: Case Studies, Tools, Dissemination and Exploitation



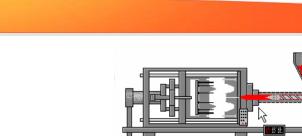
Leader: Jan Tretmans

- How to evaluate and disseminate research
 - Case Studies
 - Integration of tool components into industrial tools or tool chains
 - Continuous dissemination to industry
 - Demonstrating the applicability of our approach



Main Achievements Y1 WP5 Case Studies

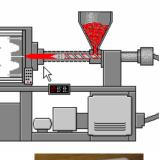
- Accumulator Charge Controller (HYDAC)
 - Simulink & Stateflow Models
 - Tool Chain: UPPAAL Tiga, Phaver, Simulink 🖡
- Wireless Sensor Network (CHESS)
 - gMAC protocol analyzed using UPPAAL (minumum waiting time for synchr.) and MODEST (prob. Of collision rates)
- Control Software for satellites Hershel and Planck (TERMA)
 - Recently started. A UPPAAL for schedulability analysis partially complete.
- Self-Balancing Scooter (CHESS)











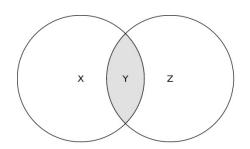






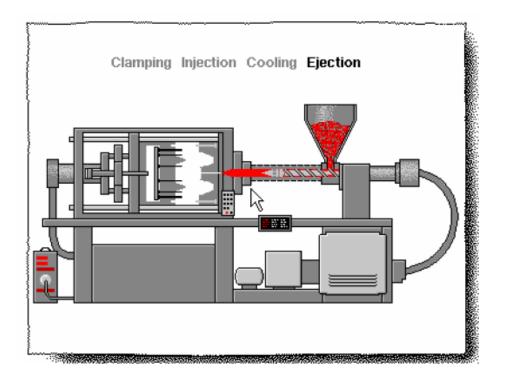
- Wireless Sensor Network (CHESS)
 - gMAC protocol analyzed using UPPAAL (minumum waiting time for synchr.) and MODEST (prob. Of collision rates)

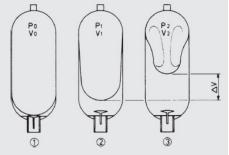


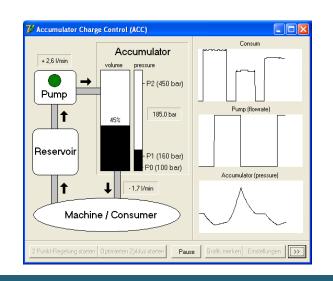




- Tool Chain: UPPAAL Tiga, Phaver, Simulink

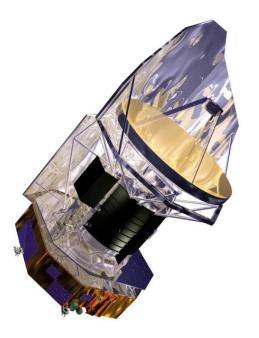








- Control Software for satellites Hershel and Planck (TERMA)
 - Recently started. A UPPAAL for schedulability analysis partially complete.







Self–Balancing Scooter (CHESS)

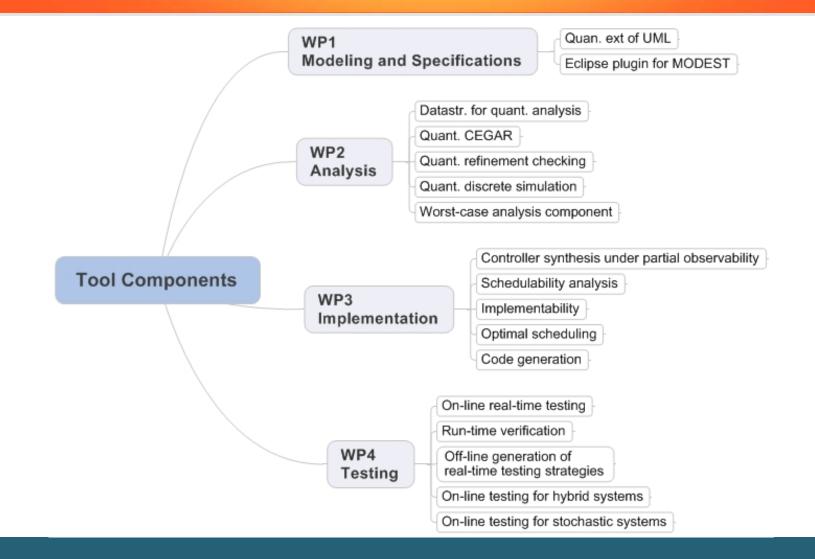


Quasimodo Review, Brussels, February 19, 2009

Page 41

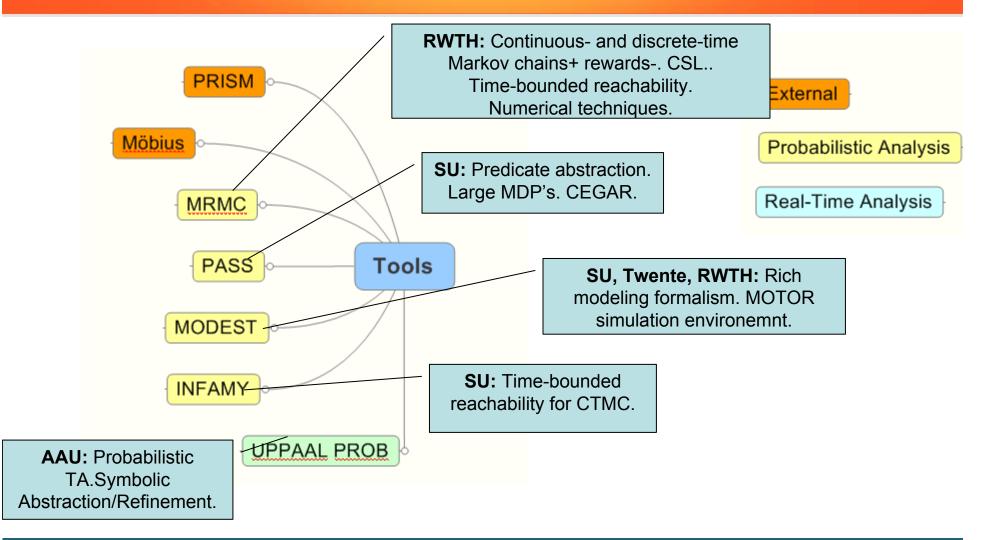
Tool Components

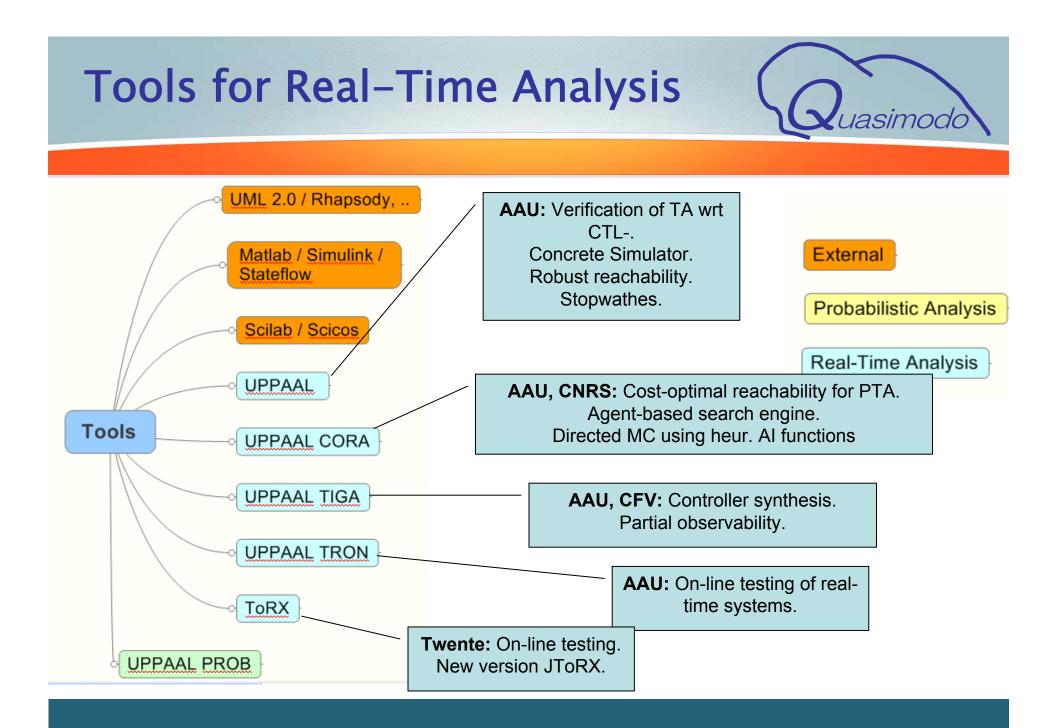




Tools for Probabilistic and Stochastic Model Checking



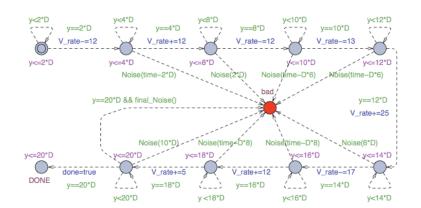


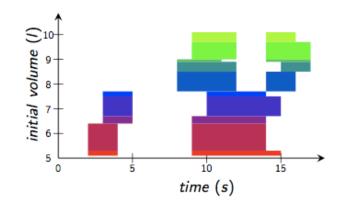


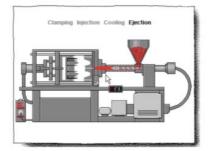
WP5-T1: Case Studies HYDAC: Accumulator Charge Controller

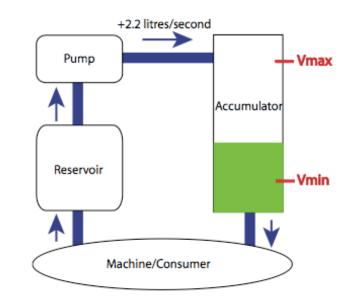


Accumulator Charge Controller (HYDAC)







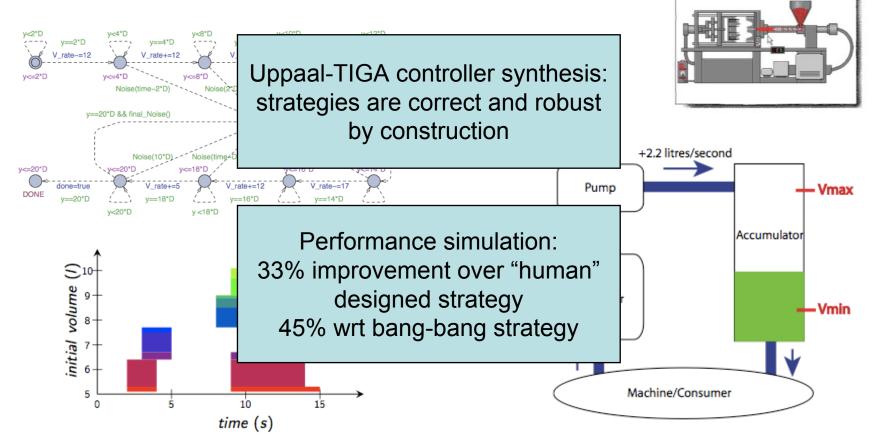


WP5-T1: Case Studies HYDAC: Accumulator Charge Controller



Clamping Injection Cooling Fiection

Accumulator Charge Controller (HYDAC)



WP5-T1: Case Studies HYDAC: Accumulator Charge Controller

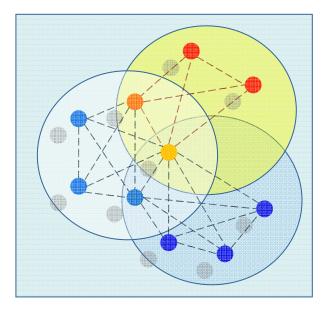


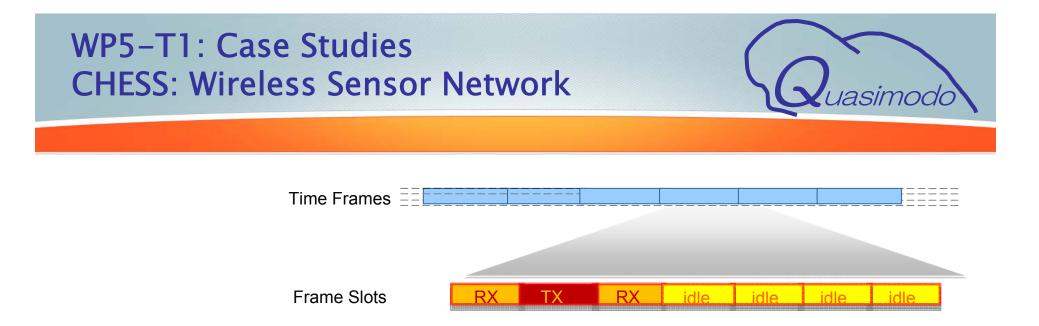
- Timed–Gamed Automata (UPPAAL–Tiga)
 - controller synthesis
 - robust (unsettled for HYDAC controllers yet)
- PHAVer
 - analysis and verification
 - HYDAC controllers are safe: pressure within margins
- Simulink and Stateflow
 - simulation, experimental validation (test of model, no proof):
 - HYDAC Smart Control uses less energy than Bang Bang Control
 - performance of synthesized controllers provide improvement
 - over HYDAC Smart Control: 33%; over Bang Bang Control: 45%



Uppaal modeling

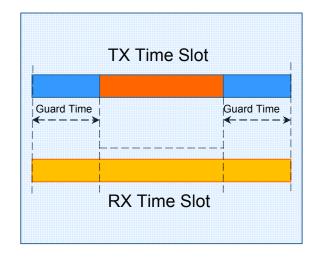
- Spatially distributed autonomous nodes
- Mobility
- Clock synchronization
 - synchronization in gMAC protocol
- Different kinds of applications



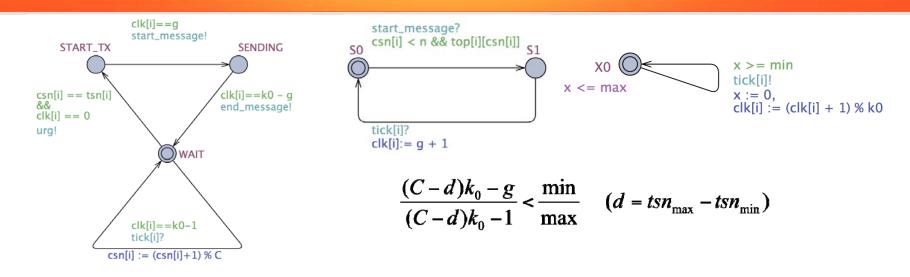


- Each node owns a slot to send (TX)
- Each node listens to other slots (RX)
- In TX slot: sender waits guard time g before and after sending a message to allow for synchronization

Which *g*? critical for functional correctness and for energy consumption







 UPPAAL models and formalization of network synchronization: if node n is sending in slot s, all other nodes are also in slot s:

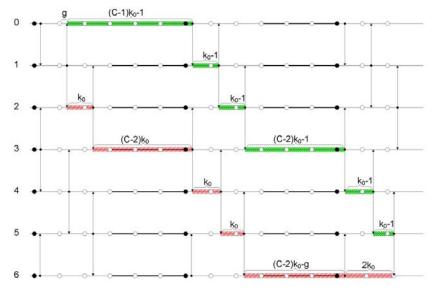
A[] forall(i:Nodes) forall(j:Nodes) (WSN(i).SENDING imply csn[i] = csn[j])

- Uppaal proves whether g leads to a synchronized network
- Generalized to formula for synchronization of connected networks



another result:

every network will eventually fail with increasing #nodes



Future

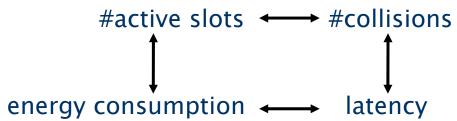
- non-fully connected networks
- different topologies
- dynamic networks in UPPAAL

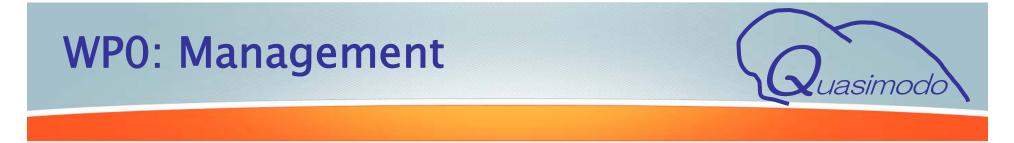
- because nodes get out of sync
- fixed guard time, #active slots, clock drift, ...
- generalizes Fan&Lynch, Meier&Thiele with decreasing clocks
- result of mathematical nature, manually proven; Uppaal used to generate counter-examples that were generalized
- practical implications?



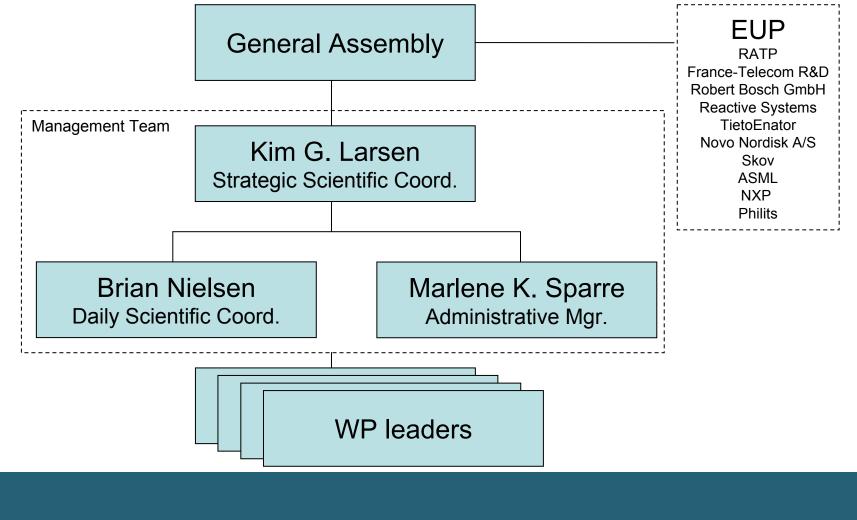
MoDeST modeling

- discrete event simulation
 - vary #active slots (now synchronization, and others, assumed)
 - determine probability of collisions vs. #active slots
 - analysis of effectiveness of collision detection mechanism
 - how this affects performance and energy consumption
- results





Efficient, flexible and reactive project management





■ 1.995 M€ EC contribution of a total 2.6 M€

~7.5 person year/year

	WP0	WP1	WP2	WP3	WP4	WP5	Total
AAU	18	2	6	6	6	8	46
ESI	0	16	11	6	14	25	72
CNRS	0	0	8	13	0	7	28
RWTH	0	4	14	0	0	10	28
SU	0	5	11	0	4	8	28
CFV	0	0	4	20	0	4	28
Terma	0	2	0	0	3	5	10
Chess	0	2	0	0	0	8	10
Inchron	0	1	0	2	2	5	10
Hydac	0	2	0	0	2	6	10
Total	18	34	54	47	31	86	270

Schedule



		_											
ID	Task Name	Otr 2	Otr 3	Qtr 4	2008	r 2 Otr	3 0#r.4	2009		3 Otr 4	2010	r2 Qtr3 Qtr4	2011
1	WP0: Management	Gauz	QUID			1 <u>1</u>			<u>202 Qu</u>				1
2	WP1: Modelling and specification				, <u> </u>	ļ.		į.	<u> </u>				ų.
3	T 1.1, D 1.3: Model process improvement	1											Ĩ.
4	T 1.2, D 1.1: Modelling quantitative system aspects	1						1					
5	T 1.3, D 1.2: Design notations	1			1	1		1	- i				i i
6	T 1.3, D1.4: Modelling tools	1											1
7	WP2: Analysis												v
8	T 2.1, D 2.1: Model checking real-time probabilistic models	1						1					
9	T 2.1, D 2.2: Symbolic data structures and analysis of multiple quantitative models	1			1			7			1		1
10	T 2.2, D 2.3: Abstraction	1						t					
11	T 2.2, D 2.4: Abstraction-refinement	1			1								1
12	T 2.3, D 2.5: Approximate analysis	1			i -	- i			i 🔤		1	i	i i
13	WP3: Implementation	1											W
14	T 3.1, D 3.1:Transfer of correctness properties from model to implementation				1	i.		Ì	- i			i i	i i
15	T 3.1, D 3.2: Tool for implementability checking	1						ł.					
16	T 3.1, D 3.5: Extended timed automata for scheduling				į.			ļ —					1
17	T 3.2, D 3.3: Model checking of controllability properties	1						1					
18	T 3.2, D 3.4: Synthesizing controllers with bounded resources	1			1	1					1		1
19	T 3.2, D 3.6: Code generation from untimed specifications	1				i		í					
20	T 3.2, D 3.7: Code generation from timed specifications	1											
21	WP4:Testing	1				i -		1	-			-i	Ŵ.
22	T 4.1, D 4.1: Quantitative testing theory	1						1					
23	T 4.1, D 4.2: Algorithms for off- and on-line quantitative testing								i.]	Ĭ		i i
24	T 4.1, D 4.3: Test selection and coverage												
25	T 4.1, D 4.5: Final algorithms and evaluation				1	1		5	1				
26	T 4.1, D 4.6: On-line hybrid/stochastic testing												
27	T4.2, D 4.4 Approximate testing					1		8					
28	WP5: Coherence, Application, Dessimination and Exploitation					i i		1	- 1				Ψ
29	T 5.1, D 5.2 Preliminary description of case studies												
30	T 5.1, D 5.5: Case studies: models				1			1					i
31	T 5.1, D 5.7: Case studies: validation										1		
32	T 5.2, D 5.4 Plan for integration of tool components				1	1		1	1			1	
33	T 5.2, D 5.8: Tool components and tool integration				1			l l					
34	T 5.2, D 5.9: Tool components				1	1		l.	1			1	1
35	T 5.1, T 5.2, D 5.10: Final report on case studies and tool integration				1								
36	T 5.3, D 5.1: Quasimodo web site					1							
37	T 5.3, D 5.3: Dissemination and use plan				i			i	- i		i i	i i	i
38	T 5.3, D 5.6: Dissemination and exploitation	na									1	- Page 54	
39	T 5.3, D 5.6: Dissemination and exploitation T 5.3, D 5.11: Final veport on dissemination and exploitation February 19, 20	03			M1	M2	M	13	M4	M	5	M69 1	VI /
40	T 5.3, D 5.12: Industrial handbook												



Quasimodo

Progress for Y1

S

Kim G. Larsen & Brian Nielsen Aalborg University, DK



- Consolidate quantitative modeling formalisms.
 - Theoretical foundation
 - Precise semantics (QLTS)
 - Formal refinement relations
- Industrial Case Studies.
 - Initial Descriptions
 - First Models.
- A first iteration of research loop:
 - Application \rightarrow Theory \rightarrow Tools \rightarrow Application





No	Deliverable name	Due Date			
D1.1	Modeling quantitative system aspects	12			
D1.2	Design notations 12				
D2.1	Model checking real-time probabilistic models	12			
D3.1	Transfer of correctness properties from model to implementation	12			
D3.3	Model checking of controllability properties	12			
D4.1	Quantitative testing theory	12			
D5.1	Quasimodo website	1			
D5.2	Preliminary description of case studies	6			
D5.3	Dissemination and use plan	6			
D5.4	Plan for integration of tool components	12			
D5.5	Case studies: models	12			

Milestones Y1



No.	Name	Date	Means of verification (Check Availability of:)		
M1	Project Start	1	Kick-off Meeting		
M2	Definition phase	6	 Precise descriptions of case studies. Plan for tool components and their integration in industrial tool chain. 		
Μ3	Modeling formalisms	12	 Semantic foundation of quantitative models in terms of labelled transition systems including semantics of composition of models, refinements between models. Formal definition of conformance and robustness between quantitative models and implementations. First models of case studies. Quantitative extensions identified by the needs of case studies. 		

Main Achievements Y1 WP1 Modeling



- Quantitative Modeling Formalisms
 - Extensive use of Timed Automata in Industrial Case Studies
 - Probabilistic Timed Automata
 - Priced Timed Automata
 - Probabilistic Priced Timed Automata
- Design Notations
 - AADL \rightarrow Probabilistic Automata
 - UML \rightarrow Timed Automata
 - UML \rightarrow Markov decision processes

Main Achievements Y1 WP2 Analysis



- Improved search engines for TA
 - Agent based | directed using Heur. Funct.
- Probabilistic Timed Automata
 - discretization | symbolic abstr./refin.
- Multi-Priced Timed Automata
- Refinement for Timed Automata as Games
- CEGAR for probabilistic programs
- Discrete-event simulation used in MRMC

Main Achievements Y1 WP3 Implementation



- Controller Synthesis
 - Fundamental results for ATL*
 - Synthesis for games with imperfect information.
 Implemented in UPPAAL-Tiga.
 - Application to HYDAC case study.
- Implementability & Code Generation
 - Robustness for Timed Automata and transfer for LTL. Implementation in UPPAAL.
 - New notion of robustness for Timed Automata based on probabilistic semantics

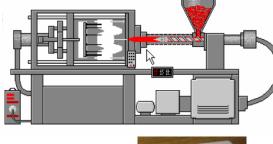
Main Achievements Y1 WP4 Testing



- Various timed ioco conformance relations defined and related.
- General quantitative ioco conformance relation provided.
- Advances on on-line testing tools for (timed) systems (ToRX, UPPAAL Tron).
- Off-line test generation as game problems using UPPAAL Tiga.

Main Achievements Y1 WP5 Case Studies

- Accumulator Charge Controller (HYDAC)
 - Simulink & Stateflow Models
 - Tool Chain: UPPAAL Tiga, Phaver, Simulink 🖡
- Wireless Sensor Network (CHESS)
 - gMAC protocol analyzed using UPPAAL (minumum waiting time for synchr.) and MODEST (prob. Of collision rates)
- Control Software for satellites Hershel and Planck (TERMA)
 - Recently started. A UPPAAL for schedulability analysis partially complete.
- Self-Balancing Scooter (CHESS)





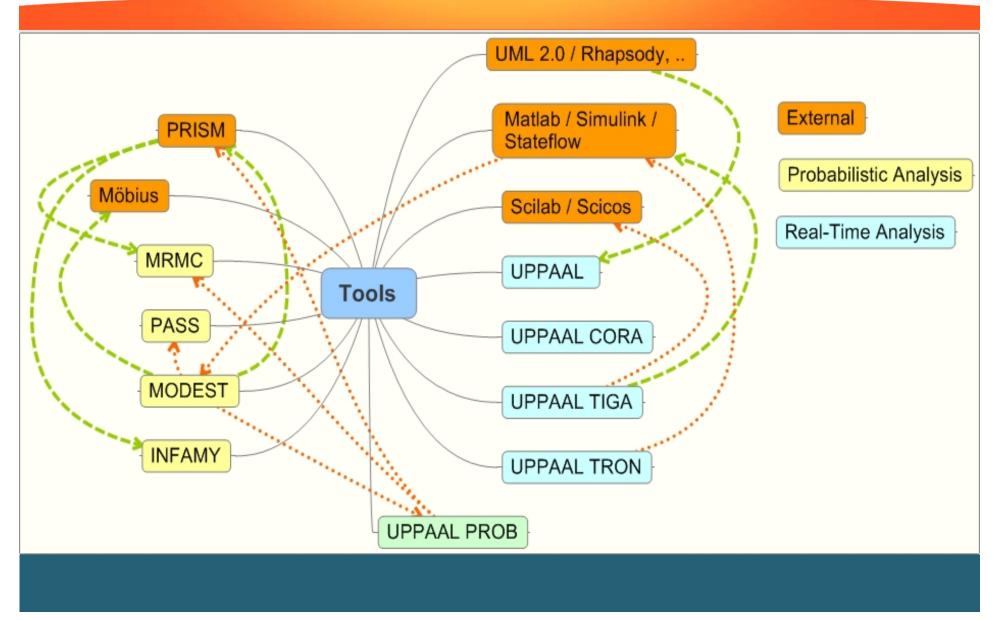






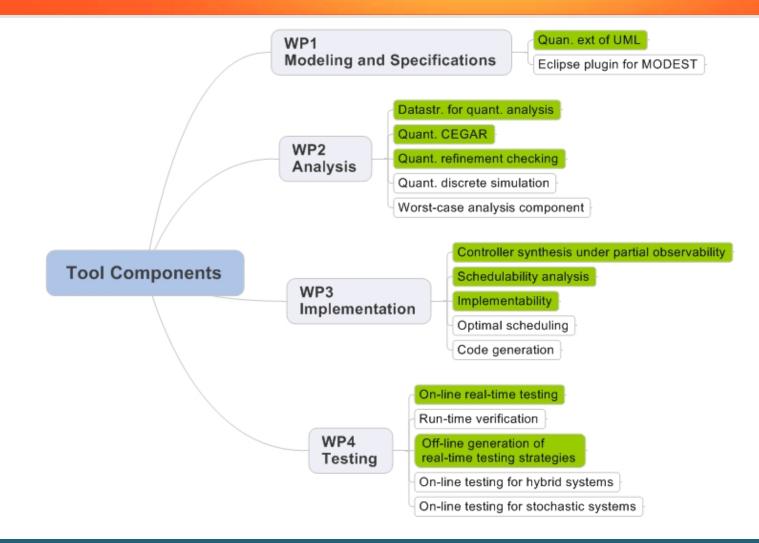
Main Achievements Y1 WP5 Tools & Tool Integration





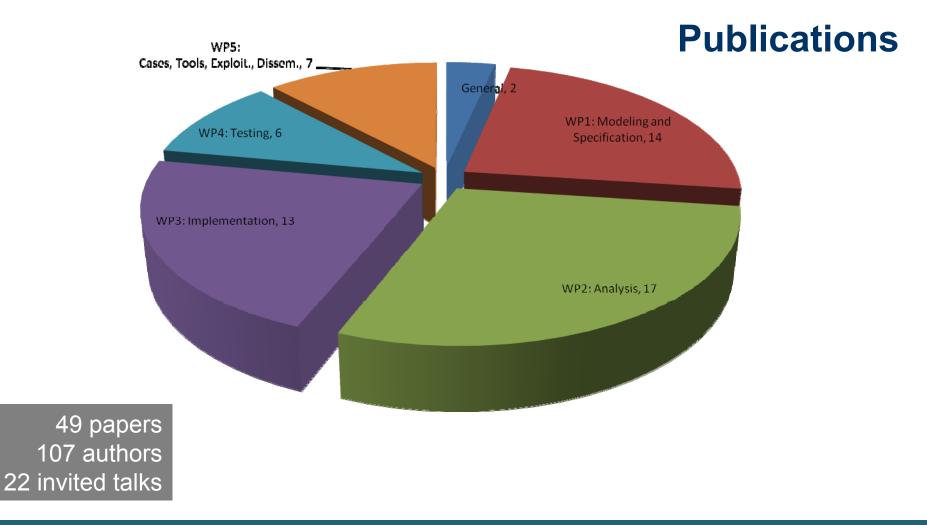
Main Achievements Y1 WP5 Tool Components





Main Achievements Y1 WP5 Dissemination





Main Achievements Y1 WP5 Dissemination



- Organisation & Contribution
 - Conferences, Workshops,
 - PhD Schools, Courses
- European Projects
 - E.g. Design, FP7 NoE
- National Projects
- Coming Events
 - FM Week Industry day at Nov. 2009, Eindhoven
 - QUANTLOG 2009 Workshop on Quantitative Logics. Rhodes, Greece, July 5–12 at ICALP 2009
 - GASICS, Workshop on Games for Design, Verification and Synthesis. at CAV 2009, Grenoble, June 26-July 2
 - Quantitative Model Checking, PhD School, Copenhagen, December 2009 (with ARTIST Design)
 - Quantitative Models: Expressivenss & Analysis, Dagstuhl Workshop, January 17–22, 2010.